A Tale of Three Signatures: Practical Attack of ECDSA with wNAF

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How to attack ECDSA

1. Focus on the primitive: DLP on elliptic curves



2. OR get extra informations from an implementation: side channel attacks.



Our work



- Improve the processing step of already known side-channel ECDSA attacks, using the Extended Hidden Number Problem and lattice techniques.
- Optimize the attack to maximize the success probability and minimize the overall time.
- Perform an attack with the minimum number of signatures needed to recover the secret key: only 3 signatures!

Our target: ECDSA

Elliptic Curve Digital Signature Algorithm is a variant of the Digital Signature Algorithm, DSA, which uses elliptic curves instead of finite fields.

Public Parameters

- An elliptic curve *E* over a prime field.
- A generator *G* of prime order *q* on *E*.
- A hash function H to \mathbb{Z}_q .

Secret Key

- An integer $\alpha \in [1, q 1]$. Public Key
 - *p_k* = [α]*G*: scalar multiplication of *G* by α.

Signing algorithm

- To sign a message *m*:
- Step 1: Randomly select nonce $k \leftarrow_R \mathbb{Z}_q$
- Step 2: Compute the point (r, y) = [k]G.
- Step 3: Compute $s = k^{-1}(H(m) + \alpha r) \mod q$.
- Step 4: Output the signature (r, s).



Step 2: Compute the point (r, y) = [k]G

- Requires a fast algorithm
- Ideally that doesn't leak any information on k!

Double-and-add algorithm

Goal: compute fast point multiplication on elliptic curves

- Input: integer k and point G.
- Output: Q = [k]G

Step 1 : Convert k to binary:

$$k = k_0 + 2k_1 + 2^2k_2 + \dots + 2^tk_t$$

Step 2 : Initialize $Q = O$
Step 3 : For $j = t, \dots, 0$, do:
• $Q \leftarrow 2Q$ double
• if $k_j = 1$: add $Q \leftarrow Q + G$
Step 4 : Return Q .

- Faster than repeated additions.
- Time of execution depends on number of 1s.
- Reduce Hamming weight of scalar k
 (w)NAF representation.

Non-adjacent form (NAF) and windowed-NAF (wNAF)

NAF:

- Impossible to have two consecutive non-zero digits,
- signed digits -1, 0, 1

wNAF:

- Impossible to have two consecutive non-zero digits,
- signed digits are in a larger window: $\in [-2^w + 1, 2^w 1]$.

Example, 3 representations of 23:

- binary: $23 = 2^4 + 2^2 + 2^1 + 2^0 = (1, 0, 1, 1, 1)$
- NAF: $23 = 2^5 2^3 2^0 = (1, 0, -1, 0, 0, -1)$
- wNAF (for w=3): $23 = 2^4 + 7 \times 2^0 = (1, 0, 0, 0, 7)$

wNAF in the wild

ECSDA with wNAF representation is used in:

- Bitcoin, as the signing algorithm for the transactions
- Some common libraries:
 - OpenSSL up to May 2019
 - Cryptlib
 - BouncyCastle
 - Apple's CommonCrypto



Oh no! Information is being leaked!

The power of side-channel attacks:

Double and add is **not** constant time (depends on the number of non-zero coeff).

 \longrightarrow (Cache) timing attacks identify (most) of the positions of the non-zero coefficients in the wNAF representation of the nonce k.

Information collected

What we have:

Many messages m_i with their signatures (s_i, r_i) , signed by a unique secret key α .



Side channels give the trace of k_i : $\star 0 0 0 \star 0 0 0 0 \star 0 0 0 \star 0 0 0 \star 0 0 0$

The important information is:

- number of non-zero coefficients, ℓ_i
- position of non-zero coefficients, $\lambda_1, \cdots, \lambda_{\ell_i}$

The Extended Hidden Number Problem

Hlavác, Rosa (SAC 2007), Extended hidden number problem and its cryptanalytic applications.

Consider u congruences of the form

$$a_i lpha + \sum_{j=1}^{\ell_i} b_{i,j} k_{i,j} \equiv c_i \pmod{q},$$

- Unknowns: the secret α and $0 \leq k_{i,j} \leq 2^{\eta_{ij}}$,
- known values: modulus $q, \eta_{ij}, a_i, b_{i,j}, c_i, \ell_i$ for $1 \leq i \leq u$,

Recover α in polynomial time.

Using EHNP to attack ECDSA

Goal: Transform ECDSA into an EHNP setup.

• ECDSA equation:

$$\alpha r = sk - H(m) \pmod{q}.$$

• Known information on the nonce k :

$$\mathbf{k} = \sum_{j=1}^{\ell} k_j 2^{\lambda_j} = \bar{k} + \sum_{j=1}^{\ell} \mathbf{d}_j 2^{\lambda_j + 1},$$

• By substitution:

$$lpha r_i - \sum_{j=1}^{\ell_i} 2^{\lambda_{i,j}+1} s_i d_{i,j} - (s_i \overline{k}_i - H(m_i)) \equiv 0 \pmod{q}$$

The Extended Hidden Number Problem

We now have u congruences of the form

$$a_i lpha + \sum_{j=1}^{\ell_i} b_{i,j} k_{i,j} \equiv c_i \pmod{q},$$

given by

$$E_i: \quad \alpha r_i - \sum_{j=1}^{\ell_i} 2^{\lambda_{i,j}+1} s_i \mathbf{d}_{i,j} - (s_i \bar{k}_i - H(m_i)) \equiv 0 \pmod{q}$$

- Unknowns: the secret key α and $0 \leq d_{i,j} \leq 2^{\mu_{i,j}}$,
- known values: modulus $q, r_i, \lambda_{i,j}, s_i, \bar{k}_i, \ell_i, H(m_i), \mu_{i,j}$ for $1 \leq i \leq u$,

Recover α in polynomial time.

HOW? \rightarrow with lattices

Reducing the size of the system

- We start with our system of modular equations E_i.
- Basic trick: Reduce the size of the system by eliminating α from the equations: $r_1 E_i r_i E_1$
 - Remember that

$$\alpha = r_1^{-1} \left(\sum_{i=1}^{\ell_1} 2^{\lambda_{1,j}+1} s_1 d_{1,j} + (s_1 \bar{k}_1 - z_1) \right) \pmod{q}.$$

• New Goal: recover the $d_{i,j}$, with a new system of equations:

$$E'_{i}: \sum_{j=1}^{\ell_{1}} \underbrace{(2^{\lambda_{1,j}+1}s_{1}r_{i})}_{:=\tau_{j,i}} \frac{d_{1,j}}{d_{1,j}} + \sum_{j=1}^{\ell_{i}} \underbrace{(-2^{\lambda_{i,j}+1}s_{i}r_{1})}_{:=\sigma_{i,j}} \frac{d_{i,j}}{d_{i,j}} - \underbrace{r_{1}(s_{i}\bar{k}_{i} - H(m_{i})) + r_{i}(s_{1}\bar{k}_{1} - H(m_{1}))}_{:=\gamma_{i}} \equiv 0 \pmod{q}.$$

Lattice: Definition, bad and good bases

Definition

A lattice is a discrete additive subgroup of \mathbb{R}^n , usually identified by a basis $\{b_1, \dots, b_n\}$.

Reduction algorithms: BKZ or LLL

Given an arbitrary basis $\{b_1, \dots, b_n\}$, find a "better" basis $\{b_1^*, \dots, b_n^*\}$.

Better \rightarrow the first vectors are shorter (and more orthogonal) in the reduced basis.



We construct a lattice such that there exists a linear combination v of the lines containing the $d_{i,j}$:

$$v = (0, \ldots, 0, d_{1,1}2^{m-\mu_{1,1}} - 2^{m-1}, \ldots, d_{u,\ell_u}2^{m-\mu_{u,\ell_u}} - 2^{m-1}, -2^{m-1}).$$

How to find v?

Goal: Find v.

- Good point: v has a particular shape
- <u>/</u>It has no reason to appear in the basis
- - 1. Make it short (by ugly manipulations of the lattice)
 - 2. Run BKZ on the basis¹
 - 3. Pray to find a good shaped vector in the reduced basis
 - 4. Try to reconstruct α with the plausible $d_{i,j}$ you get.

¹In practice $80 \leq \text{dim}(\text{lattice}) \leq 215$.

A new pre-processing method to speed-up the reduction

The slowest part of the attack: lattice reduction. BKZ reduction time \searrow if dimension \searrow OR coefficients size \searrow .

Goal: Speed up the reduction time by \searrow the size of the coefficients.

- Each trace t comes with a notion of "weight" $\mu(t)$.
- Each coefficient of the basis is multiplied by $m = \max \mu(t)$ to get integer coefficients.
- The size of the coefficients depends on *m*.

Idea: ----- pre-select traces with small weight

$$S_{a} = \{t \in \mathcal{T} | \mu(t) \leqslant a\}$$

Numerical experiment: 5000 traces from OpenSSL: $a \in [11, 67]$.

The effect of pre-processing

Key recovery time = time of 1 trial \times nbr of trials to find the key.

• Considering 4 and 5 traces with BKZ-25.



- S₁₉: already 44% of the traces
- 3 traces: from 12 days (S_{all}) to 39 h (S_{11}) on a single core.

3 ways to evaluate the attack

Several parameters need to be balanced to mount an attack:

- the preprocessing subset of traces S_a , if any
- BKZ block size β : varies between 20 and 35
 - $\beta \nearrow \Rightarrow$ probability of success of 1 trial \nearrow
 - but $\beta \nearrow \Rightarrow$ reduction time \nearrow



What is the minimal amount of signatures an attacker can use?

What are the parameters that lead to

- the fastest attack?
- the best probability of success?

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Our Main Results

- **3** signatures: 39 hours, small probability of success, *S*₁₁, BKZ-35.
- Our fastest attack:
 - 4 signatures: 1 hour 17 minutes, BKZ-25, S₁₅
 - 8 signatures: 2 minutes 25 seconds, BKZ-20, Sall
- Our most successful attack:
 - 4 signatures: 4% of success per trial, BKZ-35, Sall
 - 8 signatures: 45% of success per trial, BKZ-35, Sall

Previous attacks on ECDSA with wNAF

• **Comparing with another variant of EHNP** Fan, Wang, Cheng (CCS 2016), *Attacking OpenSSL implementation of ECDSA with a few signatures*

Attack	# signatures	Probability of success	Overall time		
[FWC2016]	5	4%	15 hours/18 minutes		
	6	35%	1 hour 21 minutes/18 minutes		
	7	68%	2 hours 23 minutes/34.5 minutes		
Our attack	3	0.2%	39 hours		
	4	4%	1 hour 17 minutes		
	5	20%	8 minutes 20 seconds		
	6	40%	5 minutes		
	7	45%	3 minutes		
	8	45%	2 minutes		

 Comparing with the Hidden Number Problem
 Van de Pol, Smart, Yarom (CT-RSA 2015) *Just a Little Bit More.* 13 signatures, 54% probability of success and 21 seconds total time to key recovery.

Errors can occur, and they often do!



Side-channel analyzis is not perfect.

Real k (wNAF) representation (unknown from an attacker):

1 0 0 0 7 0 0 0 0 0 -7 0 0 0 0 0 3 0 0 0 0 0 0 5 0 0 0 0 0 0

Information obtained by side channels:

Probability of success with various types of error

Error type 1:

A 0 coefficient misread as \star :

adds a new variable to the system, the nbr of non-zero digits is overestimated.

Error type 2:

A non-zero coefficient misread as 0: lose information necessary for key recovery.



Error 2 affects the probability of success of key recovery much more.

Resilience up to 2% of errors



- Morality: Resilience to errors up to 2% of misread digits.
- Resilience increase to 4% if we avoid certain types of errors.
- Strategy: in the side channel part, if you are not confident about your reading, choose to put a \star instead of a 0.

Thank you!

A Tale of Three Signatures: practical attack of ECDSA with wNAF Gabrielle De Micheli, Cécile Pierrot, Rémi Piau https://eprint.iacr.org/2019/861

Fastest attack

Number of	Total		Parameters	Probability of	
signatures	time	BKZ Preprocessing Δ		success (%)	
3	39 hours	35	S ₁₁	$\approx 2^3$	0.2
4	1 hour 17	25	S ₁₅	$\approx 2^3$	0.5
5	8 min 20	25	S ₁₉	$\approx 2^3$	6.5
6	3 min 55	20	S _{all}	$\approx 2^3$	7
7	2 min 43	20	S _{all}	$\approx 2^3$	17.5
8	2 min 25	20	S _{all}	$\approx 2^3$	29

Total time key recovery = time of single trial \times number of trials to find the key.

Highest probability of success of a single trial

Number of	Probability of	Parameters			Total
signatures	success (%)	BKZ	Preprocessing	Δ	time
3	0.2	35	S ₁₁	$\approx 2^3$	39 hours
4	4	35	S _{all}	$pprox 2^3$	25 hours 28
5	20	35	S _{all}	$pprox 2^3$	2 hours 42
6	40	35	S _{all}	$pprox 2^3$	1 hour 04
7	45	35	S _{all}	$pprox 2^3$	2 hours 36
8	45	35	S _{all}	$pprox 2^3$	5 hours 02

Comparing times with Fan et al, CCS 2016

Number of	Our attack	Fan et al		
signatures	Time	Success (%)	Time	Success (%)
3	39 hours	0.2%	-	-
4	1 hour 17 minutes	0.5%	41 minutes	1.5%
5	8 minutes 20 seconds	6.5%	18 minutes	1%
6	pprox 5 minutes	25%	18 minutes	22%
7	pprox 3 minutes	17.5%	34 minutes	24%
8	pprox 2 minutes	29%	-	-

Comparing success probabilities with Fan et al, CCS 2016

Number of	Our attack		Fan et al	
signatures	Success (%)	Time	Success (%)	Time
3	0.2%	39 hours	-	-
4	4%	25 hours 28 minutes	1.5%	41 minutes
5	20%	2 hours 42 minutes	4%	36 minutes
6	40%	1 hour 4 minutes	35%	1 hour 43 minutes
7	45%	2 hours 36 minutes	68%	3 hours 58 minutes
8	45%	5 hours 2 minutes	_	-

Error analysis using BKZ-25, $\Delta \approx 2^3$ and S_{all} .

Number of	Probability of success (%)					
signatures	0 errors	5 errors	10 errors	20 errors	30 errors	
4	0.28	$\ll 1$	0	0	0	
5	4.58	0.86	0.18	$\ll 1$	0	
6	19.52	5.26	1.26	0.14	$\ll 1$	
7	33.54	10.82	3.42	0.32	$\ll 1$	
8	35.14	13.26	4.70	0.58	$\ll 1$	

- Corresponds to a resilience of 2% of errors.
- Total time: 1 out of 5000 experiments, 46 sec per experiment, 65 hours on a single core